

#### 4.4 Linear programming dual

Write the dual of the following linear program

$$\begin{array}{rcll}
 \min & & 2x_2 + x_3 - 3x_4 & \\
 & x_1 - x_2 & & + 2x_4 \geq 2 \\
 & & 2x_2 + x_3 & = 4 \\
 & 2x_1 & - x_3 + x_4 & \leq 1 \\
 & & & x_1 \geq 0 \\
 & & & x_2 \geq 0 \\
 & & & x_3, x_4 \text{ unrestricted}
 \end{array}$$

#### 4.5 Dual of the transportation problem and its economic interpretation

A company produces a single type of good in  $m$  factories and stores it in  $n$  warehouses. Suppose that the factory  $i$ , with  $i \in \{1, \dots, m\}$ , has a production capacity of  $b_i$  units, and the warehouse  $j$ , with  $j \in \{1, \dots, n\}$ , requires  $d_j$  units. Let  $c_{ij}$  be the unit transportation cost from factory  $i$  to warehouse  $j$ .

Consider the problem of determining a transportation plan that minimizes the (transportation) costs, while satisfying the factory capacities and the warehouse demands.

If  $x_{ij}$  denotes the quantity of product transported from factory  $i$  to warehouse  $j$ , we have the following Linear Programming formulation:

$$\begin{array}{rcl}
 \min & \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} & \\
 & - \sum_{j=1}^n x_{ij} \geq -b_i & i \in \{1, \dots, m\} \\
 & \sum_{i=1}^m x_{ij} \geq d_j & j \in \{1, \dots, n\} \\
 & x_{ij} \geq 0 & i \in \{1, \dots, m\} \quad j \in \{1, \dots, n\}.
 \end{array}$$

Without loss of generality, we assume that the total production capacity satisfies  $\sum_{i=1}^m b_i = \sum_{j=1}^n d_j$ . Indeed, if  $\sum_{i=1}^m b_i > \sum_{j=1}^n d_j$ , we can introduce a dummy  $(n+1)$ -th warehouse with demand  $\sum_{i=1}^m b_i - \sum_{j=1}^n d_j$  and zero transportation costs.

Determine the dual of the above Linear Programming formulation and provide an economic interpretation of the dual.

